

UNITED STATES PATENT APPLICATION

FOR

Reliable Data Synchronization over Unreliable Networks

INVENTORS:

Alex Xueyuan Huang  
Luca Andrea Castellano

Prepared by:

Blakely, Sokoloff, Taylor & Zafman LLP  
12400 Wilshire Boulevard  
Seventh Floor  
Los Angeles, California 90025  
(408) 720-8300

Attorney's Docket No. 3399P046

"Express Mail" mailing label number EL617208876US

Date of Deposit July 31, 2001

I hereby certify that this paper or fee is being deposited with the United States Postal Service  
"Express Mail Post Office to Addressee" service under 37 CFR 1.10 on the date indicated above  
and is addressed to the Assistant Commissioner for Patents, Washington, D.C. 20231.

Julie Arango

(Typed or printed name of person mailing paper or fee)

Julie Arango 7/31/01  
(Signature of person mailing paper or fee)

## Reliable Data Synchronization over Unreliable Networks

[0001] This application claims the benefit of Provisional U.S. Patent application no. 60/271,081, filed on February 23, 2001, and entitled, "Method for Reliable Data Synchronization over Unreliable Networks", which is incorporated herein by reference.

### FIELD OF THE INVENTION

[0002] The present invention pertains to techniques for performing synchronization of data on multiple devices. More particularly, the present invention relates to synchronization of data on two or more devices, at least one of which is on a wireless network or other relatively unreliable network.

### BACKGROUND OF THE INVENTION

[0003] In recent years, portable electronic devices, such as laptop computers, Personal Digital Assistants (PDAs) (e.g., a Palm Pilots), and cellular telephones, have gained widespread use. These devices are frequently used, among other reasons, to carry personal information such as address books and calendars. It is not unusual for one person to own two or more such devices, often in addition to a personal computer (PC). Consequently, users of such devices generally prefer to add, delete, or modify data on one device and see the changes automatically reflected in their other devices. In response to this preference, the industry has developed data synchronization methods and standards.

[0004] Typical data synchronization methods operate based on the assumption that communication among devices is reliable, i.e., that data will not be lost

during transmission and that communication links will not fail. This is a reasonable (albeit occasionally inaccurate) assumption in the context of wired serial links and networks. Wireless links are far less reliable, however, and are likely to remain so for some time. A wireless link can fail due to many possible reasons, such as interference from nearby buildings and other structures, terrain and other electronic devices, or traveling outside service area boundaries. If synchronization fails, it may be difficult or impossible for a user to make manual adjustments to the data from a mobile device to correct for the failure.

[0005] It is expected that within a short time, wireless links will be the predominant connection medium between personal electronic devices. It is crucial, therefore, that these devices be capable of synchronizing their data without errors, even when communication links fail or devices on the network are offline.

## SUMMARY OF THE INVENTION

[0006] The present invention includes a method and apparatus to synchronize states of data between a number of devices over an unreliable communication channel. In the method, data is retrieved from the devices. Based on the data retrieved from the devices, centrally stored data is updated so as to automatically recover from a prior synchronization failure, if any. The data states on the devices are then updated based on the updated centrally stored data, including communicating with at least one of the devices over the unreliable communication channel.

[0007] Other features of the present invention will be apparent from the accompanying drawings and from the detailed description which follows.

## BRIEF DESCRIPTION OF THE DRAWINGS

[0008] The present invention is illustrated by way of example and not limitation in the figures of the accompanying drawings, in which like references indicate similar elements and in which:

[0009] Figure 1 illustrates a network environment in which the present invention can operate;

[0010] Figure 2 illustrates a physical implementation of the environment in which the invention operates;

[0011] Figure 3 shows the elements of the sync manager, according to one embodiment;

[0012] Figure 4 shows the state of each of multiple devices and the action table in an initial state;

[0013] Figure 5 shows the state of each device and the action table after the cell phone has been successfully updated;

[0014] Figure 6 shows the state of the system after synchronization is complete;

[0015] Figure 7 is a flow diagram illustrating an overall process for assuring reliable synchronization over an unreliable network, according to one embodiment;

[0016] Figure 8 is a flow diagram illustrating in greater detail the operation of determining actual contact states, according to one embodiment;

[0017] Figure 9 is a flow diagram illustrating in greater detail the operation of producing a new truth database and action table, according to one embodiment;

and

[0018] Figure 10 is a high-level block diagram of a processing system.

Figure 10 is a high-level block diagram of a processing system.

## DETAILED DESCRIPTION

[0019] A method and apparatus for fault-tolerant synchronization of data between two or more devices, at least one of which is on a relatively unreliable network, are described. Note that in this description, references to “one embodiment” or “an embodiment” mean that the feature being referred to is included in at least one embodiment of the present invention. Further, separate references to “one embodiment” in this description do not necessarily refer to the same embodiment; however, neither are such embodiments mutually exclusive, unless so stated and except as will be readily apparent to those skilled in the art. For example, a feature, structure, act, etc. described in one embodiment may also be included in other embodiments. Thus, the present invention can include a variety of combinations and/or integrations of the embodiments described herein.

[0020] As used herein, an “unreliable” network is a network which is appreciably more subject to communication failures (even though such failures may have short duration) than a conventional, wired communications network. One example of an unreliable network is a conventional wireless telecommunications network, such as used for cellular telephony. Note, however, that the fault-tolerant synchronization techniques described herein are not limited to use with a wireless network.

[0021] The techniques described herein assure that data will be accurately synchronized among two or more devices, even if a communication link fails or a

device is switched off during synchronization. As described in greater detail below, this goal is accomplished by operating a centralized synchronization system which maintains a "truth database" indicating the true state of the data and "action tables" indicating actions to be performed on each of the devices the next time the data is updated. During a synchronization process, the synchronization system first retrieves the data to be synchronized from each of the online devices, via the network, and determines actual states of the data on the devices based on the action tables and the retrieved data. The synchronization system updates the truth database and action tables accordingly. The synchronization system then generates an effective action table for any of the devices which were offline during the previous synchronization. The truth database and effective action tables are then stored in an atomic transaction. The synchronization system then use the effective action tables to update the data on the devices over the network. As will be more apparent from the details which follow, this process provides automatic recovery from synchronization failures, such as may result from failure of communication links or devices being offline during synchronization.

[0022] Figure 1 shows an example of a network environment in which this technique may be implemented. A centralized synchronization ("sync") system 1 synchronizes a set of data between a number of remote sync points 2. In one embodiment, synchronization is performed using SyncML. A "sync point" is a user device containing data to be synchronized. The sync system 1 may be



implemented in one or more physical computing platforms, which may be one or more conventional server-class computers and/or personal computers (PCs).

Some of the sync points 2 are connected to the sync system 1 via the Internet 3.

However, one or more other sync points 2 are connected to the sync system 1 via a wireless network 4, which may be a conventional wireless telecommunications network. The sync points 2 on the Internet 3 may include devices such as conventional desktop and laptop PCs, while the sync points 2 on the wireless network 4 may include devices such as cellular telephones, PDAs, laptop computers, and two-way pagers.

[0023] The sync system 1 includes a sync manager 5 and a data storage server (DSS) 6 coupled to the sync manager 5. The sync manager 5 coordinates the synchronization process. The sync manager 5 also generates a conventional user interface 7, to allow users of sync points on the Internet to add, delete, and modify their data and to initiate synchronization. The DSS 6 stores and retrieves binary large objects and (BLOBs) "atomically"; that is, the DSS 6 guarantees that when an element is written or updated, it will either update completely, or it will not update at all. In an embodiment which handles personal information manager (PIM) data (i.e., personal contact data), one BLOB is created for each user, containing all of the contact information belonging to that user. To facilitate description, it is henceforth assumed in this description that the data to be synchronized is PIM data. It will be recognized, however, that the techniques described herein can be used to synchronize essentially any type of data between

multiple devices.

[0024] The sync manager 5 constructs a truth database 8 reflecting the sync manager's notion of what the true, synchronized state of all the sync points 2 should be, based on the information available to the sync manager 5 from the sync points 2. The sync manager 5 also maintains an action table 9 for each sync point 2, to record the changes that need to be made to each sync point 2 in order to make the data on the sync points 2 consistent with the truth database. When the data on all of the sync points 2 are consistent with the data in the truth database, synchronization is complete.

[0025] Each entry in an action table contains an operator, and (in some cases) associated data. The operator is one of: NO\_CHANGE, ADD, DELETE or UPDATE. Action table entries using the ADD and UPDATE operators also have data associated with the operator.

[0026] To facilitate implementation, the action tables for the various sync points 2 can be separate only in a logical sense; that is, they may be defined logically within a single, larger action table, if desired. If so, each entry in the action table also contains a sync point identifier (SPID), which uniquely identifies the sync point to which that entry applies.

[0027] Figure 2 illustrates an example of the physical implementation of the environment in which the invention operates. In the embodiment of Figure 2, the sync system 1 includes the sync manager 5, the DSS 6, and a conventional World Wide Web server 21 (hereinafter simply "Web server 21"), all coupled to

each other on a LAN 22. The user interface 7 in Figure 1 may take the form of a conventional Web page, which is provided to a sync point 2 by the Web Server 21. The Web Server 21 is connected to a sync point 2 (a PC in this case,) via the Internet 3.

[0028] The sync manager 5 communicates with at least one mobile synch point 2 (e.g., a cellular telephone) via the wireless network 4. Although not shown as such in Figure 2, the sync manager 5 may be coupled to the wireless network 4 via the Internet 3 or some other wired network. In that case, communication between the sync manager 5 and the wireless network 4 may be carried out through a gateway (not shown) that provides an interface between the wireless 4 network and the Internet 3. Such a gateway may operate to allow wireless devices to access the Web, email, and other services that require Internet access. Hence, such a gateway may perform functions such as translation of data between hypertext markup language in (HTML) and wireless markup language in (WML) and translating protocols between hypertext transfer and protocol (HTTP) and wireless accessed protocol (WAP), or other similar conversions. Such a gateway can be implemented using, for example, the Mobile Access Gateway software, available from Openwave Systems of Redwood City, California, with conventional server-class computer hardware.

[0029] As shown in Figure 2, the DSS 6 may include a contact application program interface (API) 23, a BLOB server 24, and a data storage facility 25. The BLOB server 24 is a conventional database server that provides access to contact

records stored in the data storage facility 25. The data storage facility 25 is any conventional device (or multiple devices) suitable for the non-volatile mass storage of data. The contact API 23 is an API which allows communication of personal contact data (names, addresses, telephone numbers, etc.) between the BLOB server 24 and elements external to the DSS 6. The BLOB server 24 allows atomic reads and writes of BLOBs to the data storage facility 25. The details of the contact API 23 and the BLOB server 24 are not necessary for an understanding of the present invention and are well within the abilities of those skilled in the art, given this disclosure.

[0030] Figure 3 shows the elements of the sync manager 5, according to one embodiment. As shown, the sync manager 5 includes a sync engine 31, a recovery module 32, and a preferences module 33. The sync engine 31 is a conventional data synchronization engine, which is responsible for propagating changes between sync points, resolving conflicting states between two or more sync points, and detecting duplicate records. The recovery module 32 provides the fault-tolerance features described herein. The preferences module 33 allows user preferences to be set (e.g., prioritization of sync points for purposes of conflict resolution). In one embodiment, the sync engine 31, recovery module 32 and preferences module 33 are embodied as separate software-implemented processes of the sync manager 5.

[0031] Figures 4, 5 and 6 collectively illustrate a simple example of the use of the synchronization technique described herein, in which no failure occurs.

Reference numeral 40 indicates the data states on each of the three sync points 2. To facilitate description, a single action table 9 is used for all sync points 2 in this example, rather than separate action tables for each sync point 2. Each entry in the action table 9 is associated with a particular sync point 2 using a four-digit SPID.

[0032] Figure 4 shows the initial PIM state of three sync points 2, i.e., a cell phone, a notebook computer, and a desktop computer, and the action table 9. Figure 5 shows the PIM state of each sync point 2 and the action table 9 after the data on the cell phone has been successfully updated (i.e., to update the telephone number of Mitsuko Uchida), based on the second entry in the action table 9. Figure 6 shows the state of the system after synchronization is complete. Each sync point 2 has the same contact data and the action table 9 is empty.

[0033] In this example, every component and communication link worked as expected. That is, all of the communication links delivered their data without errors, and every sync point and network element stayed online and worked correctly. In practice, however, failures will sometimes occur during synchronization. Some of the possible failure modes are: the machine running the sync manager 5 fails, the DSS 6 fails, the network connection between the sync manager 5 and the DSS 6 fails, or the network connection with a sync point 2 (wired or wireless) fails during synchronization.

[0034] Typical synchronization methods do not necessarily recover from these failures. Failure of current systems to recover from network and device lapses

can result in the loss and/or corruption of user data. The result is often dissatisfied users. The technique described herein assures that user data will not be corrupted even if such failures occur.

[0035] The technique can be viewed as having three main aspects. First, the truth database is built from three elements: the truth database from the previous synchronization, the current sync point state, and any incomplete actions from previous synchronizations. This aspect of the technique provides the ability to recover from a failed synchronization.

[0036] The second aspect is that, after constructing the truth database and the action table, but before updating any of the sync points, the sync manager 5 writes both the truth database and the action table to the DSS 6 in an atomic transaction. When the sync manager 5 receives an acknowledgement from a sync point 2 that that device has successfully synchronized, the sync manager 5 removes the corresponding entries from the action table and writes the revised truth database and action table back to the DSS 6. This aspect of the technique ensures that proper recovery information is always available.

[0037] The third aspect is that, during a synchronization, when a sync point 2 is offline and there is an old action table for that sync point 2, the old action table is combined with the new action table to form an effective action table to be used during the next synchronization.

[0038] Figure 7 illustrates a top-level process to assure reliable synchronization, according to one embodiment of the present invention. The illustrated sequence

is performed at each synchronization. Synchronization may be either user-initiated (e.g., by the user activating a predetermined control on a sync point) or server-initiated. Note that a two-stage process is employed to assure persistence of data, even in the event of a server failure: first, information is gathered from all sync points 2 to determine the truth database and action table; then, data is written to the DSS 6 and to the sync points 2. If a failure occurs before the data can be written to the DSS 6, the sync manager 5 may simply declare all data modified since the last synchronization to be invalid and reinitiate the process.

[0039] Initially, at processing block 701 the sync manager 5 reads from the DSS 6, in an atomic transaction, the truth database and any action tables remaining from a previous synchronization. At block 702 the sync engine 31 retrieves all contacts and their states from each of the sync points. At block 703 the recovery module 32 determines, for each contact of each sync point 2, the actual current state of the contact data on the sync point. This determination is made based on both the action table for the sync point and the state of the contact as retrieved from the sync point, as described further below. The sync engine 31 then produces a new truth database and action tables for each of the sync points 2 at block 704. At block 705, the sync engine 31 produces an effective action table for each sync point that was offline (if any) during the last synchronization. Any effective action table is produced by merging the previous action table with the new action table for each such sync point 2, as described further below. Next, at block 706 the sync engine 31 writes the truth database and the effective action

tables to the DSS 6 in anatomic transaction. For purposes of this write operation, the effective action table of any sync point which was online during the previous synchronization is simply the new action table produced in block 704. At block 707 the sync engine 31 updates each sync point 2 in succession, via the appropriate network, using well-known, conventional techniques, in which each step of the update requires acknowledgement from the sync point being updated. Finally, at block 708 the sync engine 31 deletes the action table of each sync point 2 for which the update was successful.

[0040] Figure 8 shows in greater detail the process of block 703 in which the recovery module 32 determines the actual states of the data records on a sync point 2, according to one embodiment. This process is carried out for each individual data element of each data record on each sync point 2. Assuming an embodiment which synchronizes PIM data, for example, each record may represent a different personal contact and may include multiple individual data elements, such as address, telephone number, e-mail address, etc. Note that there are four possible data states: ADDED, UPDATED, DELETED, NO\_CHANGE. Note also the difference between a data state and a "to do" action: Data states are in the past tense (e.g., ADDED, UPDATED, DELETED), while "to do" actions are in the present tense (e.g., ADD, UPDATE, DELETE).

[0041] There are two determining factors in recovery: 1) what the action table recorded as the action to do to a record, and 2) the state returned from the sync point for that record. Thus, the process illustrated in Figure 8 uses these two



factors. Initially, at block 801 it is determined whether the action in the action table for this sync point is DELETE. If so, then this record (e.g., a personal contact) is simply deleted at block 802; that is, it is not presented to the sync engine 31 for use in generating a new truth database and action tables (block 704 of the top-level process). In this case, the record is not presented to the sync engine 31, because on the last synchronization, the sync engine 31 must have already deleted it from the truth database. Note that the "yes" branch from block 801 can apply in any of four possible cases, i.e., when the state of the handset is DELETED, ADDED, NO\_CHANGE or UPDATED.

[0042] If the action is determined not to be DELETE in block 801, then at block 803 it is determined whether the state of the record, as indicated by the sync point 2, is DELETED. If the state is DELETED, then at block 804 any actions for this record are removed from the action table of this sync point, and the state of this record on the sync point is deemed to be DELETED at block 805. Note that if the record is deleted from the sync point, there is no need for any of the actions from the action table relating to this record, because they do not apply.

Therefore, they can be removed.

[0043] If instead the sync point state is determined not to be DELETED at block 803, then it is determined at block 806 whether the state of the record from the sync point 2 is UPDATED. If the state is UPDATED, then at block 807 any actions for this record are removed from the action table of this sync point 2, and the state of this record on the sync point 2 is deemed to be DELETED at block

808. Note that in this case, if the action in the action table is UPDATE for this record and sync point, then this simply means that the sync point has been updated by the user, such that whatever is in the action table for this record is irrelevant. If the action is ADD, it means the communication link was broken after the sync point added the record but before the sync point sent an acknowledgement; now the user has updated the record on the sync point 2. If the action is NO\_CHANGE, the record was updated on the sync point 2. If the action is DELETE, this case is handled by block 801.

[0044] If the sync point state is determined not to be UPDATED at block 803, then it is determined at block 809 whether the state of the record from the sync point 2 is ADDED. If not, the process ends with block 820, in which the state of this record on the sync point is deemed to be NO\_CHANGE. In that case, the action in the action table is simply maintained, as if the sync point was offline. If, however, the sync point state is found to be ADDED at block 809, the process continues from block 810. At block 810 any conventional duplicate detection technique is used to determine whether the record from the sync point 2 is a duplicate of another record. This case might occur if the record added by the sync point 2 is actually an existing record which the sync point 2 mistakenly interpreted as new. If the record is not a duplicate, the process ends with block 821, in which the state of this record on the sync point 2 is deemed to be ADDED. If the sync point record is a duplicate of another record (the "old" record), then the process continues from block 811, in which it is determined whether the

action table of this sync point 2 specifies ADD for the old record. If not, the process continues from block 815. If the action table specifies ADD for the old record, the process continues from block 812.

[0045] If the action table does not specify ADD at block 811, then at block 815 it is determined whether the action table specifies UPDATE for the old record. If so, then the process ends with block 816, in which the record from the sync point 2 is deleted (i.e., not presented to the sync engine). Note that this is an unlikely scenario. If it occurs, it probably means that the sync engine 31 determined that the record is supposed to be on the sync point 2 and attempted to update it; however, the sync point 2 believes the record has been added.

[0046] If the action table does not specify UPDATE for the old record at block 815, it is most likely that this is because the user added this record to the sync point 2 himself. In that event, at block 817 the new record from the sync point 2 is assigned the identifier of the old record, after which the old record is modified as necessary to make it match the new record at block 818. The process then ends.

[0047] If the action table is found at block 811 to specify ADD for the old record, then at block 812 the action for the old record is removed from the action table. Following block 812, the record from the sync point 2 is deleted at block 813 (i.e., not presented to the sync engine 31). In that case, the process ends with block 814, in which the "new" record from the sync point 2 is assigned the identifier of the old record. This case can occur when the sync engine 31

requested the sync point 2 to add this entry on the last synchronization, and it was added, but an acknowledgement was not received from the sync point 2.

[0048] Figure 9 shows in greater detail the process of producing a new truth database and action tables (block 704 in Figure 7), according to one embodiment. Initially, at block 901 the sync engine 31 examines a record. At block 902 the sync engine 31 determines whether there is only a single change in state for this record, i.e., whether only one sync point 2 has changed the state of the record. If there is only a single change in state, the process branches to block 905, in which the sync engine 31 modifies the truth database to reflect the change. Following block 905, the sync engine 31 records the intent to propagate the change in the action tables of all other sync points 2 at block 912. If there are additional records to examine (block 910), the sync engine 31 accesses the next record at block 911, and the process then loops back to block 901.

[0049] If there is not only one change in the state of the record at block 902, then at block 903 the sync engine 31 determines whether there is a change in the state of the record from multiple sync points 2. If so, then at block 904 the sync engine 31 uses any suitable conflict resolution technique to choose from which sync point 2 the state of the record should be taken, for purposes of updating the truth database. As one example of conflict resolution, the sync engine 31 may simply use a previously set user preference indicating the user's preferred prioritization of his sync points 2. Following block 904, the process continues from block 905 as described above.

[0050] If there was not a change in state from multiple sync points 2 at block 903, then it is determined at block 906 whether the state of the record is ADDED (as determined by the process of Figure 8) . If so, then at block 907 the sync engine 31 determines whether the added record is a duplicate of an old record. If so, the process continues from block 909, in which the sync engine 31 simply ignores the record. If the added record is not a duplicate, then at block 908 sync engine 31 adds the record to the truth database, and the process continues from block 912 as described above. If the sync engine 31 determines the state of the record is not ADDED at block 906, then the sync engine 31 ignores the record at block 909. Following block 909, the process continues from block 910. If there are no more records to examine at block 910, the process ends.

[0051] As noted above, the third main aspect of the synchronization technique is the method by which the sync manager 5 handles an offline sync point 2. The sync manager 5 combines the action table constructed from the online sync points and the truth database (the "new" action table) with the action table residing in the DSS 6 from the prior synchronization (the "old" action table) to produce an effective action table. The effective action table ensures that the appropriate synchronization actions will take place the next time the sync point 2 is online during synchronization.

[0052] The process for producing the effective action table is as follows, according to one embodiment. The new action table and the old action table are collectively referred to as the "input" action tables. First, all of the action table

entries which do not appear in both input action tables are entered directly into the effective action table. For each remaining entry, the input operators are mapped to an output operator using the table below. In the table below, the action from the new action table selects the column, while the corresponding action from the old action table selects the row. The selected entry is the action that is placed in the corresponding entry of the effective action table. (Note that the asymmetry in the table is an effect of the passage of time.) If the output operator has data associated with it (as with ADD and MODIFY actions), the corresponding data from the truth database data is used.

Operation	No Change	Add	Delete	Update
No Change	No Change	Add	Delete	Update
Add	Add	Add	No Change	Add
Delete	Delete	Invalid	Delete	Invalid
Update	Update	Update	Delete	Update

[0053] For example, suppose the offline sync point 2 had an incomplete action to ADD a certain record. Assume further that subsequently, another sync point 2 UPDATED that record. The UPDATE column and the ADD row yields an output action of ADD, to be placed in the effective action table for that record. This means that at the next synchronization, the offline sync point 2 will have this record added to its database, but with the updated contents.

[0054] Figure 10 is a block diagram of a processing system that may represent

any of the communication and/or processing devices mentioned above (e.g., sync system 1 or sync points 2.). Note that the sync system 1 in particular may be distributed between two or more processing systems such as shown in Figure 10. The illustrated processing system includes one or more processors 101, i.e. a central processing unit (CPU), read-only memory (ROM) 102, and random access memory (RAM) 103, each connected to a bus system 107. The processor(s) 101 may be, or may include, one or more programmable general-purpose or special-purpose microprocessors, digital signal processors (DSPs), application specific integrated circuits (ASICs), programmable logic devices (PLDs), or a combination of such devices. The bus system 107 includes one or more buses or other connections, which may be connected to each other through various bridges, controllers and/or adapters, such as are well-known in the art. For example, the bus system 107 may include a "system bus", which may be connected through one or more adapters to one or more expansion buses, such as a Peripheral Component Interconnect (PCI) bus, HyperTransport or industry standard architecture (ISA) bus, small computer system interface (SCSI) bus, universal serial bus (USB), or Institute of Electrical and Electronics Engineers (IEEE) standard 1394 bus (sometimes referred to as "Firewire").

[0055] Also coupled to the bus system 107 are one or more mass storage devices 104, input/output (I/O) devices 105 and data communication devices 106. Mass storage device 104 may be, or may include, any one or more devices suitable for storing large volumes of data in a non-volatile manner, such as a

magnetic disk or tape, magneto-optical (MO) storage device, or any of various forms of Digital Versatile Disk (DVD) or Compact Disk (CD) based storage, or a combination thereof.

[0056] Each data communication device 106 is a device suitable for enabling the processing system to communicate data with remote devices or systems and may be, for example, a conventional modem, a Digital Subscriber Line (DSL) modem, a cable modem, an Ethernet adapter, an Integrated Services Digital Network (ISDN) adapter, a wireless transceiver (e.g., in the case of a mobile device), a satellite transceiver, or the like. The I/O device(s) 105 may include, for example, a keyboard or keypad, a display device, and a pointing device (e.g., a mouse, trackball, or touchpad). Note, however, that such I/O devices may be unnecessary for certain devices and/or in certain embodiments. For example, a device which functions purely as a server does not necessarily require local I/O devices aside from a data communication device, particularly if the server is not intended to directly interface with a user or operator. Similarly, it may not be necessary or desirable to equip a mobile sync point with a mass storage device.

Many other variations on the above described embodiment other possible.

Further, it will be understood that the processing system may include other conventional components such as are well-known in the art.

[0057] The sync processes described above may be implemented in software 108, which may reside, either partially or completely, in any of RAM 103, mass storage device 104 and/or ROM 102, as shown.



1203 1204 1205 1206 1207 1208 1209 1210 1211 1212 1213 1214 1215 1216 1217 1218 1219 1220 1221 1222 1223 1224 1225 1226 1227 1228 1229 1230 1231 1232 1233 1234 1235 1236 1237 1238 1239 1240 1241 1242 1243 1244 1245 1246 1247 1248 1249 1250 1251 1252 1253 1254 1255 1256 1257 1258 1259 1260 1261 1262 1263 1264 1265 1266 1267 1268 1269 1270 1271 1272 1273 1274 1275 1276 1277 1278 1279 1280 1281 1282 1283 1284 1285 1286 1287 1288 1289 1290 1291 1292 1293 1294 1295 1296 1297 1298 1299 1300 1301 1302 1303 1304 1305 1306 1307 1308 1309 1310 1311 1312 1313 1314 1315 1316 1317 1318 1319 1320 1321 1322 1323 1324 1325 1326 1327 1328 1329 1330 1331 1332 1333 1334 1335 1336 1337 1338 1339 1340 1341 1342 1343 1344 1345 1346 1347 1348 1349 1350 1351 1352 1353 1354 1355 1356 1357 1358 1359 1360 1361 1362 1363 1364 1365 1366 1367 1368 1369 1370 1371 1372 1373 1374 1375 1376 1377 1378 1379 1380 1381 1382 1383 1384 1385 1386 1387 1388 1389 1390 1391 1392 1393 1394 1395 1396 1397 1398 1399 1400 1401 1402 1403 1404 1405 1406 1407 1408 1409 1410 1411 1412 1413 1414 1415 1416 1417 1418 1419 1420 1421 1422 1423 1424 1425 1426 1427 1428 1429 1430 1431 1432 1433 1434 1435 1436 1437 1438 1439 1440 1441 1442 1443 1444 1445 1446 1447 1448 1449 1450 1451 1452 1453 1454 1455 1456 1457 1458 1459 1460 1461 1462 1463 1464 1465 1466 1467 1468 1469 1470 1471 1472 1473 1474 1475 1476 1477 1478 1479 1480 1481 1482 1483 1484 1485 1486 1487 1488 1489 1490 1491 1492 1493 1494 1495 1496 1497 1498 1499 1500 1501 1502 1503 1504 1505 1506 1507 1508 1509 1510 1511 1512 1513 1514 1515 1516 1517 1518 1519 1520 1521 1522 1523 1524 1525 1526 1527 1528 1529 1530 1531 1532 1533 1534 1535 1536 1537 1538 1539 1540 1541 1542 1543 1544 1545 1546 1547 1548 1549 1550 1551 1552 1553 1554 1555 1556 1557 1558 1559 1560 1561 1562 1563 1564 1565 1566 1567 1568 1569 1570 1571 1572 1573 1574 1575 1576 1577 1578 1579 1580 1581 1582 1583 1584 1585 1586 1587 1588 1589 1590 1591 1592 1593 1594 1595 1596 1597 1598 1599 1600 1601 1602 1603 1604 1605 1606 1607 1608 1609 1610 1611 1612 1613 1614 1615 1616 1617 1618 1619 1620 1621 1622 1623 1624 1625 1626 1627 1628 1629 1630 1631 1632 1633 1634 1635 1636 1637 1638 1639 1640 1641 1642 1643 1644 1645 1646 1647 1648 1649 1650 1651 1652 1653 1654 1655 1656 1657 1658 1659 1660 1661 1662 1663 1664 1665 1666 1667 1668 1669 1670 1671 1672 1673 1674 1675 1676 1677 1678 1679 1680 1681 1682 1683 1684 1685 1686 1687 1688 1689 1690 1691 1692 1693 1694 1695 1696 1697 1698 1699 1700 1701 1702 1703 1704 1705 1706 1707 1708 1709 1710 1711 1712 1713 1714 1715 1716 1717 1718 1719 1720 1721 1722 1723 1724 1725 1726 1727 1728 1729 1730 1731 1732 1733 1734 1735 1736 1737 1738 1739 1740 1741 1742 1743 1744 1745 1746 1747 1748 1749 1750 1751 1752 1753 1754 1755 1756 1757 1758 1759 1760 1761 1762 1763 1764 1765 1766 1767 1768 1769 1770 1771 1772 1773 1774 1775 1776 1777 1778 1779 1780 1781 1782 1783 1784 1785 1786 1787 1788 1789 1790 1791 1792 1793 1794 1795 1796 1797 1798 1799 1800 1801 1802 1803 1804 1805 1806 1807 1808 1809 1810 1811 1812 1813 1814 1815 1816 1817 1818 1819 1820 1821 1822 1823 1824 1825 1826 1827 1828 1829 1830 1831 1832 1833 1834 1835 1836 1837 1838 1839 1840 1841 1842 1843 1844 1845 1846 1847 1848 1849 1850 1851 1852 1853 1854 1855 1856 1857 1858 1859 1860 1861 1862 1863 1864 1865 1866 1867 1868 1869 1870 1871 1872 1873 1874 1875 1876 1877 1878 1879 1880 1881 1882 1883 1884 1885 1886 1887 1888 1889 1890 1891 1892 1893 1894 1895 1896 1897 1898 1899 1900 1901 1902 1903 1904 1905 1906 1907 1908 1909 1910 1911 1912 1913 1914 1915 1916 1917 1918 1919 1920 1921 1922 1923 1924 1925 1926 1927 1928 1929 1930 1931 1932 1933 1934 1935 1936 1937 1938 1939 1940 1941 1942 1943 1944 1945 1946 1947 1948 1949 1950 1951 1952 1953 1954 1955 1956 1957 1958 1959 1960 1961 1962 1963 1964 1965 1966 1967 1968 1969 1970 1971 1972 1973 1974 1975 1976 1977 1978 1979 1980 1981 1982 1983 1984 1985 1986 1987 1988 1989 1990 1991 1992 1993 1994 1995 1996 1997 1998 1999 2000 2001 2002 2003 2004 2005 2006 2007 2008 2009 2010 2011 2012 2013 2014 2015 2016 2017 2018 2019 2020 2